







# The Impact of Kernel Asynchronous APIs on the Performance of a Kernel VPN

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### Context and Motivation

Virtual Private Networks (VPNS): For privacy and security



\$68.3 Billion Market Share (2025)<sup>1</sup>



More than 1.75 Billions of users<sup>2</sup>

WireGuard: modern, fast VPN in the Linux Kernel



Is Wireguard, a multi-threaded VPN kernel module able to handle thousands of clients efficiently?

<sup>&</sup>lt;sup>1</sup>https://www.coherentmarketinsights.com/industry-reports/virtual-private-network-market

<sup>&</sup>lt;sup>2</sup>https://surfshark.com/blog/vpn-users

### Question: Does WireGuard scale?

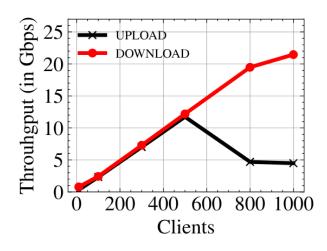
#### Evaluation:

- 0 1,000 clients \* 25Mbps = 25 Gbps generated traffic
- Use cases: Client Upload, Client Download
- 25 Gbps Mellanox Connect-X 4 NIC
- 18 Cores Intel Xeon Gold 5220

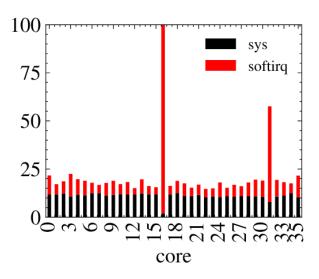
#### Metrics

- The forwarded network throughput by the server
- The Server CPU Usage

### Problem: WireGuard doesn't scale!



Reception Throughput



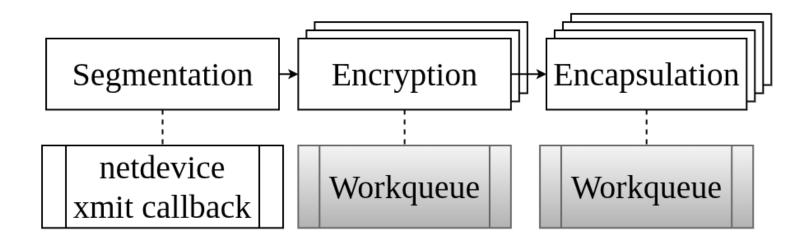
CPU usage

- Download use case scales well
- Upload use case doesn't scale!
  - Peaks at 12 Gbps with 500 clients
  - Plateaus at 4.5 Gbps (500+ clients)
- CPU usage is not 100% at 1,000 clients
  - o CPU is not the bottleneck
  - NIC is not the bottleneck

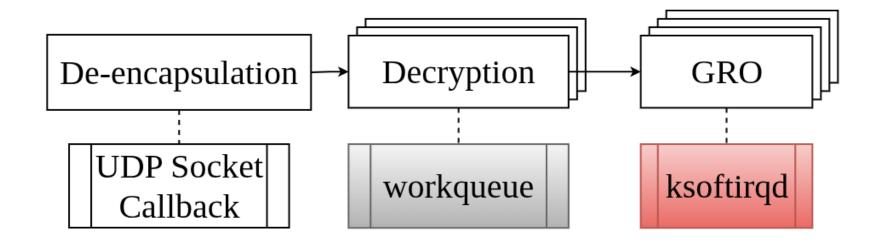
## WireGuard Linux Kernel Implementation

- Performs:
  - Traffic Encryption
  - UDP Tunnelling
- Implemented on top of Linux network stack
  - Leverages Generic Segmentation/Receive Offload (GSO/GRO)
- Uses Linux kernel asynchronous APIs:
  - Workqueues
  - Softirq
- Two pipelines:
  - Transmission pipeline
  - Reception pipeline

### Transmission Pipeline

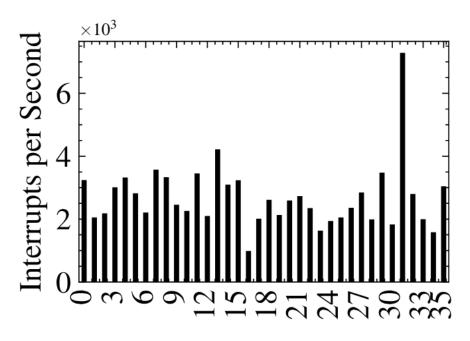


## Reception Pipeline



### Bottleneck is not trivial!

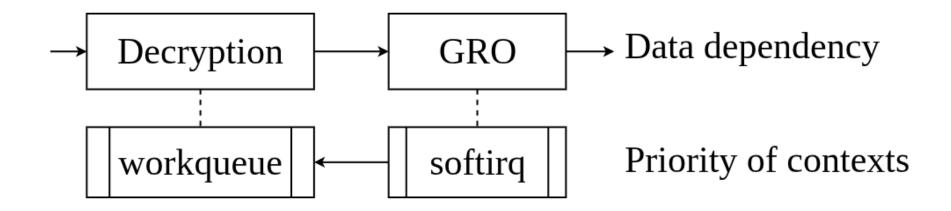
- Checked multiple possible reasons:
  - Network packet distribution across CPUs (RSS) works fine
  - Workqueues are fine, as there is no problem in the Download use case where transmission pipeline is sollicited
  - WireGuard is multi-threaded



Net RX Interrupts distribution

So what is the real problem?

## The Bottleneck is Execution Order Inversion



Definition: later pipeline stages preempt earlier ones due to priority mismatch

## Impact of Eol

Ordered (Optimal): Jobs Processed = 2, Latency = 3

F1 F2 F3 F1 F2 F3 Context G1

1 job processed 1 job processed Context G2

Unordered: Jobs Processed = 1, Latency = 3+ε

F1 F3 F2 F3 F1 F3 Context G3

T job processed

- EoI happens in 80% of all the jobs 'processing
- Eol increases latency and decreases throughput

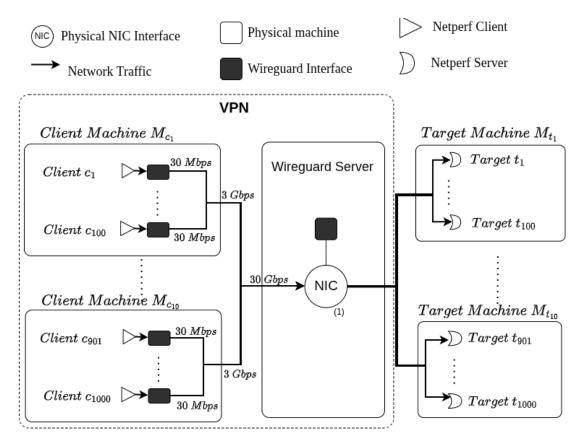
# Solution: Using Different Asynchronous APIs

- Run GRO at the same priority as decryption to preserve order
- Two alternatives for GRO execution:
  - Kthreads (threaded NAPI)
  - Workqueues (new extension to NAPI)
- Kthreads
  - Easy to deploy (config only)
  - But one thread per client -> scalability issues
- Workqueues
  - Requires kernel + Wireguard changes
  - Fixed thread pool (per CPU) -> Scalable

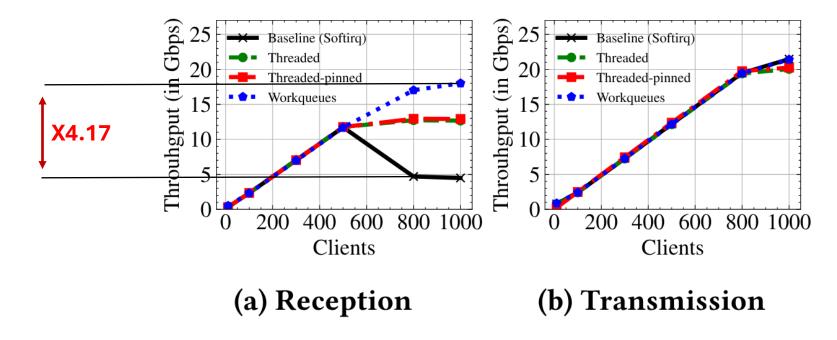
### **Evaluation and Results**

#### Evaluation Setup:

- o Testbed: 21 servers
- Intel Xeon Gold 5220 (18 cores)
- 25 Gbps NIC, Mellanox Connect-X 4
- Linux 6.1 (LTS), Debian 12
- Evaluation Scenario
  - Up to 1,000 clients generating each
     25Mbps of traffic upload and download with iPerf3
- Metrics:
  - Throughput, CPU Usage, 99th Tail Latency

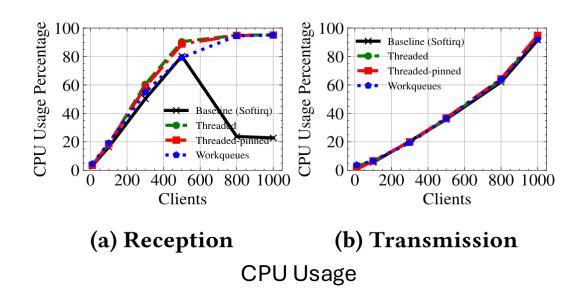


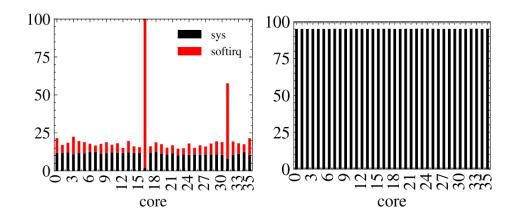
## Evaluation and Results: Throughput



- In reception, with 800+ clients, throuhgput:
  - o remains at 12.5 Gbps with kthreads (a x2.8 Improvement)
  - Scales with workqueueus up to 18.8 Gbps
    - A x4.17 improvement
- Transmission pipeline is not impacted, which is good.

# Evaluation and Results: CPU Usage





(a) Baseline (softirq) (b) kthreads/workqueues
Per core CPU usage

CPU is now fully used

# Evaluation and Results: Latency

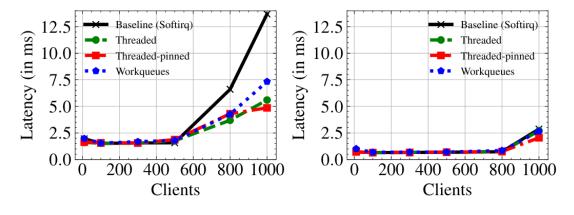
#### • Upload:

- o Baseline 13.7 ms
- o kthreads 5.6 ms (4.8 ms pinned, best)
- o workqueue 7.3 ms.

#### Download:

- Low latency overall (0.6–0.7 ms up to 500 clients)
- rises at scale (baseline 2.8 ms, kthreads/workqueue 2.7 ms, pinned 2.0 ms, best)

Pinned kthreads consistently achieve the lowest latency, especially under high load.



(a) Client upload (b) Client download 99th Tail Latency

### Takeaways

When designing a multi-threaded asynchronous application, the choice of which execution context to use is crucial, even more so for kernel modules.